### **Comparative Cheap Talk**

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#### Cheap talk about private information

- Seller knows something about quality of a product
- Professor knows something about prospects of a student
- Analyst knows something about value of a stock
- Auditor knows something about viability of a company
- Ebay knows something about quality of sellers

#### Interval cheap talk (Crawford and Sobel, 1982)

- Biased Sender S knows realization of r.v. θ distributed uniformly on [0,1]
- Receiver R takes action a

 $U^{S} = -(a-(\theta+b))^{2}$  $U^{R} = -(a-\theta)^{2}$ 

- Informed R chooses  $a = \theta$
- But S ideal is  $a = \theta + b$
- For any realization  $\theta$  that S reports, R will believe true  $\theta$  is lower
- So S must exaggerate even more and R compensates even more
- So communication breaks down and R just chooses  $a=E[\theta]$
- S receives expected payoff of only  $E[U^S] = -Var[\theta]-b^2$
- Which is less than if S told truth!  $E[U^S] = -b^2$

# But some communication is still possible for small b

- Consider b=1/10
- S states θ in [0,3/10] or θ in (3/10,1]
- S ideal action for  $\theta = 3/10$  is a = 4/10
- E[θ| θ in [0,3/10]]= 3/20
- E[θ| θ in (3/10,1]]=13/20
- So S is indifferent at  $\theta$ =3/10 and prefers truth on either side
- Smaller b, more partitions b>1/4 no communication b<1/4 two partitions max b<1/12 three partitions max etc.

## If interests are too far apart such cheap talk does not work

- Uncertainty over quality of a good seller tells buyer it's great
- Uncertainty over prospects for a stock analyst says it's a sure bet
- Uncertainty over whether a spending program is worthwhile – administrator says it's essential
- Uncertainty over quality of a job applicant recommender says he's great

Cheap talk is not credible if there is some action the sender always wants the receiver to take (e.g. the maximal action)

## What if there is uncertainty along multiple dimensions?

- Seller has multiple goods they're all great!
- Analyst recommends multiple stocks buy them all!
- Lobbyist favors multiple spending programs they're all necessary!
- Professor has multiple students they will all win Nobel prizes!

# Can comparative statements be credible?

- Good A is better than B
- Stock A is better than B
- Proposal A is better than B
- Student A is better than B
- Comparative statements are positive along one dimension and negative along another dimension at the same time
- Can't exaggerate!
- But still might have an incentive to invert the ordering

### The symmetric model

•  $\theta = (\theta_1, \dots, \theta_N)$ 

(density strictly positive on  $[0,1]^N$  and different  $\theta_k$  i.i.d.)

- a=(a<sub>1</sub>,..., a<sub>N</sub>) (each action in [0,1], can choose actions independently)
- Payoffs additive across dimensions:  $U^{R}(a,\theta) = \Sigma_{k} u^{R}(\theta_{k},a_{k})$

 $U^{S}(a,\theta) = \Sigma_{k} u^{S}(\theta_{k},a_{k})$ 

- Complete ordering: rank variables from worst to best
- Partial ordering: categorize variables into groups from worst to best (and don't differentiate within a group)

#### **Recommendation game**

Professor (S) has N students to recommend to employer (R). Each student hired (a<sub>k</sub> = 1) or not (a<sub>k</sub> = 0).

 $u^{S} = (\theta_{k} - T^{S})a_{k}$  $u^{R} = (\theta_{k} - T^{R})a_{k}$ 

- Employer hires a student if expected quality  $\theta_k$  above T<sup>R</sup>
- Professor wants a student to be hired if  $\theta_k$  above  $T^S$
- "CS" equilibrium in one dimension if  $E[\theta_k|\theta_k < T^S] < T^R < E[\theta_k|\theta_k > T^S]$

# Recommendation game – comparative cheap talk

- Two students k and k', where  $\theta_k > \theta_{k'}$
- $E[\theta_{j:N}] = j/(N+1)$
- Assume T<sup>R</sup>=3/5 so top student hired, bottom student not hired
- Sender payoff from correct ordering:  $(\theta_k T^S)1 + (\theta_{k'} T^S)0$
- Sender payoff from inverted ordering:  $(\theta_k T^S)0 + (\theta_{k'} T^S)1$
- So gain from deviation:  $\theta_{k'} \theta_k < 0$
- True even if T<sup>S</sup>=0 so professor always wants a student to be hired so CS interval cheap talk impossible
- Both employer and professor are happier if a better student is hired than if a worse student is hired – both utility functions supermodular

## Same idea can be applied to standard uniform-quadratic C-S game

For all b, sender and receiver payoffs are supermodular ("sorting condition"), so any ordering is an equilibrium.

$$u^{R} = -(a_{k} - \theta_{k})^{2} \qquad u_{12}^{R} = 2 > 0$$
  

$$u^{S} = -(a_{k} - (\theta_{k} + b))^{2} \qquad u_{12}^{S} = 2 > 0$$
  

$$a_{j:N} = E[\theta_{j:N}] = j/(N+1)$$

So for any two j> j', and  $\theta_k > \theta_{k'}$  sender won't deviate if:  $-(E[\theta_{j':N}] - (\theta'+b))^2 - (E[\theta_{j:N}] - (\theta+b))^2 \ge$   $-(E[\theta_{j':N}] - (\theta+b))^2 - (E[\theta_{j:N}] - (\theta'+b))^2$ or  $(E[\theta_{j:N}] - E[\theta_{j':N}])(\theta_k - \theta_{k'}) \ge 0$ 

- Theorem 1: If  $u^R$  and  $u^S$  are supermodular in  $(\theta_k, a_k)$  then the complete ordering and every partial ordering are equilibrium orderings.
- S states ordering, worst to best
- R chooses action  $a_{i:N}$  for jth worst issue
- Sender won't lie if for any j>j' and  $\theta_k > \theta_{k'}$ :

 $u^{S}(\theta, a_{j:N}) + u^{S}(\theta', a_{j':N}) \geq u^{S}(\theta', a_{j:N}) + u^{S}(\theta, a_{j':N})$ 

- But if  $a_{i:N} > a_{i':N}$  this is just supermodularity
- So when does receiver take higher action for higher ranked issue?
- Receiver takes action to maximize utility. If knew θ for certain then supermodularity good enough – but only knows ranking of θ.
- Ranking implies  $\theta_{j:N} > _{FOSD} \theta_{j':N}$ .
- FOSD plus supermodularity implies that action is higher for higher ranked variables

## Rankings become more informative as number of issues increases



Distribution of top-ranked issue

- Theorem 2: Under the complete ordering, expected sender and receiver payoffs asymptotically approach the full information case as the number of issues N increases.
- Application to uniform-quadratic game:
- Sender payoffs are concave:

 $\begin{aligned} \mathsf{u}^{\mathsf{R}} &= -(\mathsf{a}_{\mathsf{k}}\text{-}\boldsymbol{\theta}_{\mathsf{k}})^2 \\ \mathsf{u}^{\mathsf{S}} &= -(\mathsf{a}_{\mathsf{k}}\text{-}(\boldsymbol{\theta}_{\mathsf{k}}\text{+}\mathsf{b}))^2 \end{aligned}$ 

- Babbling per-issue payoff: E[u<sup>R</sup>] = -Var[θ<sub>k</sub>] E[u<sup>S</sup>] = -Var[θ<sub>k</sub>]-b<sup>2</sup>
- Complete ordering payoff for issue j: E[u<sup>R</sup>] = -Var[θ<sub>j:N</sub>] E[u<sup>S</sup>] = -Var[θ<sub>j:N</sub>]-b<sup>2</sup>
- In complete ordering  $Var[\theta_{i:N}]$  goes to 0 in limit:

 $E[u^{R}] = 0$  $E[u^{S}] = -b^{2}$ 

#### Effect on expected sender payoffs



For given N, partition cheap talk outperforms ordinal cheap talk for small enough bias b

For given b, ordinal cheap talk outperforms partition cheap talk for large enough N

# What if asymmetric preferences, distributions?

- Theorem 3 (roughly): For N=2 comparative cheap talk is generically robust to small asymmetries.
  - Have to take rankings with "a grain of salt" and adjust for sender bias
  - Large enough asymmetries then cheap talk can break down
  - But in some cases still works for arbitrary asymmetries in preferences
  - Information loss from asymmetries
- Theorem 4 (roughly): For large enough N there are always some issues similar enough to be compared.

As before let  $a_{j:N}$  be R's best action for issue ranked *j*th from the bottom. And let  $a(\theta_k)$  be R's best action when  $\theta_k$  is known. For each q (0,1), by the Glivenko-Cantelli Theorem

 $\lim_{N\to\infty} \theta_{[qN]:N} = F^{-1}(q)$  a.s.

where  $F(\theta_k)$  is distribution of each  $\theta_k$ .

Therefore for i=S,R,  $\lim_{N\to\infty} (1/N) \sum_{i} E[u^{i}(a_{j:N},\theta_{j:N})] = E[u^{i}(a(\theta_{k}),\theta_{k})]$ 

So in the limit as the number of issues increases, the sender reveals **all** information and sender and receiver payoffs are equivalent

## Is complete order always better than partial order?

Recommendation game as before,  $T^R=3/5$ ,  $T^S=0$ Babbling/No Information:

•  $E[\theta_k]=1/2$  so no one is hired

Complete ordering:

- N=2: E[θ<sub>1:2</sub>]=1/3, E[θ<sub>2:2</sub>]=2/3
- N=3:  $E[\theta_{1:3}]=1/4$ ,  $E[\theta_{2:3}]=2/4$ ,  $E[\theta_{3:3}]=3/4$
- N=4:  $E[\theta_{1:4}]=1/5$ ,  $E[\theta_{2:4}]=2/5$ ,  $E[\theta_{3:4}]=3/5$ ,  $E[\theta_{4:4}]=4/5$
- N=5:  $E[\theta_{1:5}]=1/6$ ,  $E[\theta_{2:5}]=2/6$ ,  $E[\theta_{3:5}]=3/6$ ,  $E[\theta_{4:5}]=4/6$ ,  $E[\theta_{5:5}]=5/6$

Partial ordering:

- N=3:  $E[\theta_{1:3}]=1/4$ ,  $E[\theta_{\{2,3\}:3}]=5/8$ ,
- N=4:  $E[\theta_{\{1,2\}:4}]=3/10$ ,  $E[\theta_{\{3,4\}:4}]=7/10$
- N=5:  $E[\theta_{\{1,2\}:5}]=1/4$ ,  $E[\theta_{\{3,4,5\}:5}]=2/3$

### Conclusion

- Multiple dimensions increase the scope for communication
- Simple rankings are often credible when other forms of cheap talk are not
- These rankings can be surprisingly informative
- Sometimes sender prefers a partial ordering
- And sometimes only a partial ordering is credible
- Asymmetries, private receiver info reduce scope for communication but often still possible

### Extensions

- Interdependent actions
  - Can only hire one person
  - Must buy goods in bundle
- Sender and receiver take actions, e.g. bargaining
- Private receiver information, e.g. auctions
- Non-additive payoffs
- Reputation